Minimum Spanning Tree Optimizations

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CS 4299: Undergraduate Research
Presentation Structure

- Introduction
- Other people’s approaches
  - Classical approaches
  - Modern approaches
- Our approaches
- Result analysis
- Conclusion
Introduction

- The Minimum Spanning Tree (MST) problem is a graph problem.
- The input graph has the following properties.
  - It must be a single connected component.
  - It must be an undirected graph.
  - It must have edge weights.
- Here are the properties of the output graph.
  - It must be a tree (no cycles).
  - It must contain a connection to all nodes from the initial graph.
  - It must be the lightest possible graph that follows the previous conditions.
Introduction

- This problem has many important applications
  - Network design
    - Communication networks
    - Road networks
    - Utility networks
  - Approximation for NP-hard problems
    - Traveling Salesman Problem
    - Steiner Trees
  - Computer Science
    - Computer vision
    - Handwriting recognition
    - Image registration
    - Cluster analysis
  - Other various applications
    - Biology
    - Medical
    - Engineering
    - Finance

- Computer Networks
- Road Networks
- Medical Maps
- TSP Approximation
- Biology
For Prim’s Algorithm you do the following:

1. Add any node to the MST
2. Find the lightest edge leading outside of the MST
3. Add the destination node to the MST
4. Repeat until you have an MST
Classical Approaches: Kruskal’s

- For Kruskal’s Algorithm you do the following
  1. Sort edge list
  2. Check if the lightest edge creates a cycle
  3. If it does not create a cycle add it
  4. Repeat until you have an MST
For Boruvka’s Algorithm you do the following

1. First make every node a tree
2. Find the lightest edge leading outside of the tree
3. Merge trees together that are connected via these lightest edges
4. Repeat until you have an MST
Modern Approaches

• Prim
  - Priority queues
    • Binary
    • Binomial
    • Fibonacci
  - Fredman and Tarjan

• Kruskal
  - Bucket sorting
  - Filter Kruskal

• Boruvka
  - Yao
  - MST verification
  - Combinations
  - Other various
    • Non-Greedy
    • Vertex removal
Our Approaches: Linked List Removal

- Most graphs have a large number of nodes that only connect to 1 edge
  - These edges have to be in the MST
    - Remember the spanning property
- If this edge leads to a node with 2 edges
  - These edges have to also be in the MST
  - This is due to the fact that it takes 2 edges that could form a cycle
    - We know that an edge leading to single node cannot form a cycle
    - So you can add the other edge that this node is connected to
- This leads to improvement
  - You do not have to track as many trees
  - Fewer nodes to sort (Kruskal)
Our Approaches: Skipping Edges

- With Boruvka you can determine that certain edges cannot be in the MST
  - The edge creates a cycle
  - The non-lightest edge connecting the same trees
- You can also determine the same thing about nodes as well
  - If there are no viable edges
Input Graphs: Edge & Nodes

Edges & Nodes of Input Graphs

- internet
- USA-road-d-NV
- USA-road-d-BAY
- USA-road-d-DY
- USA-road-d-FLA
- USA-road-d-NW
- USA-road-d-NE
- 2d-2e20.sym
- USA-road-d-CAL
- USA-road-d-LYS
- USA-road-d-E
- USA-road-d-W
- USA-road-d-CTR
- USA-road-d-USA
- 4e-2e23.sym

Edges & Nodes: 11
Edges Per Node Distribution For All Graphs

- 17% of graphs have 1 edge per node
- 26% have 2 edges per node
- 45% have 5 to 8 edges per node
- 7% have 9 to 16 edges per node
- 5% have 17 to 32 edges per node
- 17% have 33 to 64 edges per node
- 5% have 65 to 128 edges per node
- 5% have 129 to 256 edges per node

Input Graphs: Degree Distribution
ECL Codes Comparison

- **Serial Version**
  - Boruvka
    - Baseline
    - Edge Skipping
    - Linked List Removal
  - Kruskal
    - Baseline
    - Bucket Sorting
    - Linked List Removal
  - Prim
    - Edge based
    - Node based

- **OpenMP**
  - Boruvka
  - Kruskal
  - Prim
Result Analysis: Comparison Code

• Serial
  • Boost
    • Industry standard
    • Well documented
    • General graph solving C++ library
    • Kruskal and Prim version
Result Analysis: Serial Runtime Comparison

On average our Kruskal baseline code was 4.5 times faster than Boost’s.

On average our Boruvka baseline code was 1.07 times faster than Boost.
This led to a 1.5 times improvement over Boost.
This led to a 5.5 times improvement over Boost
Result Analysis: OpenMP Runtime Comparison

On average our Kruskal OMP was 2.3 times faster than our serial

On average our Boruvka OMP code was 3 times faster than our serial
Conclusion

- Our code out performs Boost
  - Our Boruvka is 1.5 times faster
  - Our Kruskal is 5.5 times faster
- Be careful when relying on big O
- Don’t forget that smaller optimizations can still lead to improvements
  - Edge Skipping
  - Linked List Removal
- Sometimes it is better to break your problem into individual parts if you can
  - For example: The connectivity and the weights can be broken into separate problems
Questions?