

Assertions

Reading assignment

- A. J. Offutt, "A Practical System for Mutation Testing: Help for the Common Programmer," Proceedings of the 12th International Conference on Testing Computer Software, Washington, D.C., June 1995, pp. 99-109.

Reading assignment

- L. A. Clarke, A. Podgurski, D. J. Richardson and Steven J. Zeil, "A Formal Evaluation of Data Flow Path Selection Criteria," *IEEE Transactions on Software Engineering*, 15 (11), November 1989, pp. 1318-1332.
- Background reading
 - S. Rapps and E. J. Weyuker, "Data Flow Analysis Techniques for Test Data Selection," *Proceedings of the Sixth International Conference of Software Engineering*, Tokyo, Japan, September 1982, pp. 272-277.

Assertions

- Self-checking software
- insert specifications about the intent of a system
 - Violation means there is a **fault** in the system
- during execution, monitor if the assertion is violated
- if violated then report the violation

History of Assertions

- Alan Turing discussed using “assert statements” in algorithms, ~1947
- Assert statements used in formal verification to indicate what should be true at points in a program, ~1967
- Assertions advocated for finding faults during execution, ~1972
 - Based on preprocessors

History of Assertions

- Assertions introduced as part of programming and specification languages, 1975- >
 - Euclid, Alphard, Clu, ...
- Bertrand Meyer popularizes Design by Contract and includes assertions as an integral part of Eiffel, an OO language
- Assertion capabilities for common programming languages, available but limited

Parts of an Assertion Mechanism

- a high-level language
 - for representing logical expressions (typically Boolean-valued expressions) for characterizing invalid program execution states
 - for associating the logical expressions with well-defined states of the program (scope of applicability)
- automatic translation of the logical expressions into executable statements that evaluate the expressions on the appropriate states of the associated program
- predefined or user-defined runtime response that is invoked if the logical expression is violated

Language for representing logical expressions

- Usually use a notation that can be “easily” translated into the programming language
- Boolean expressions
 - Use variables and operators defined in the program
 - Must adhere to programming languages scoping rules
 - `ASSERT X < Y + Z;`

where X, Y, and Z are variables in the program
- Quantification
 - ForAll and ThereExists

Example

- `--ASSERT` for all I , $(1 \leq I < N)$,
 $A[I] \leq A[I + 1]$
- `--ASSERT` for some I , $(1 \leq I < N)$,
 $A[I] \leq A[I + 1]$
- not always supported since quantification can result in expensive computation

Language for representing logical expressions

- Want to reference original value and current value of a variable
 - $\text{Pre}(X)$ versus X
 - $\text{Old}(X)$ versus X
 - X' versus X

Example old and current values

- `--ASSERT` for all I , ($1 \leq I \leq N$),
 $\text{old}(A[I]) = A[I]$
 - Value of the array has not changed
- `--ASSERT` for all J , ($1 \leq J \leq N$)
(for some I , ($1 \leq I \leq N$), $\text{old}(A[J]) = A[I]$)
 - Permutation of the array

Scope of an assertion

- Local assertion
 - checked at the definition site
- Global assertion
 - defined over a specific scope, usually using the scoping rules of the programming language
 - must determine the locations that need to be checked,
 - Global ASSERT $X > 10$
must determine all the locations where X is defined and check that X is greater than 10

Scope of an assertion

- **Loop assertion (Loop invariant)**
 - Checked at each iteration at the designated point in a loop
- **Class assertion (Class invariant)**
 - Checked at the start and end of each method in a class
- **Pre (and Post conditions)**
 - Checked at the start (and end) of a method each time it is invoked
- **All of the above are syntactic sugar**

More Advanced Assertion Language Capabilities

- may be able to introduce additional (hidden) operators, operands, and types
 - e.g., length operator for stack
 - must be able to define the hidden entities in terms of the provided entities
 - `--ASSERT Z < Bound (Q)`
 - means that whenever Z is assigned a value, it must be less than Bound (Q),
 - where Bound(Q) is visible wherever Z is visible and
 - either Bound(Q) is already defined in the program or is defined to be a hidden operation

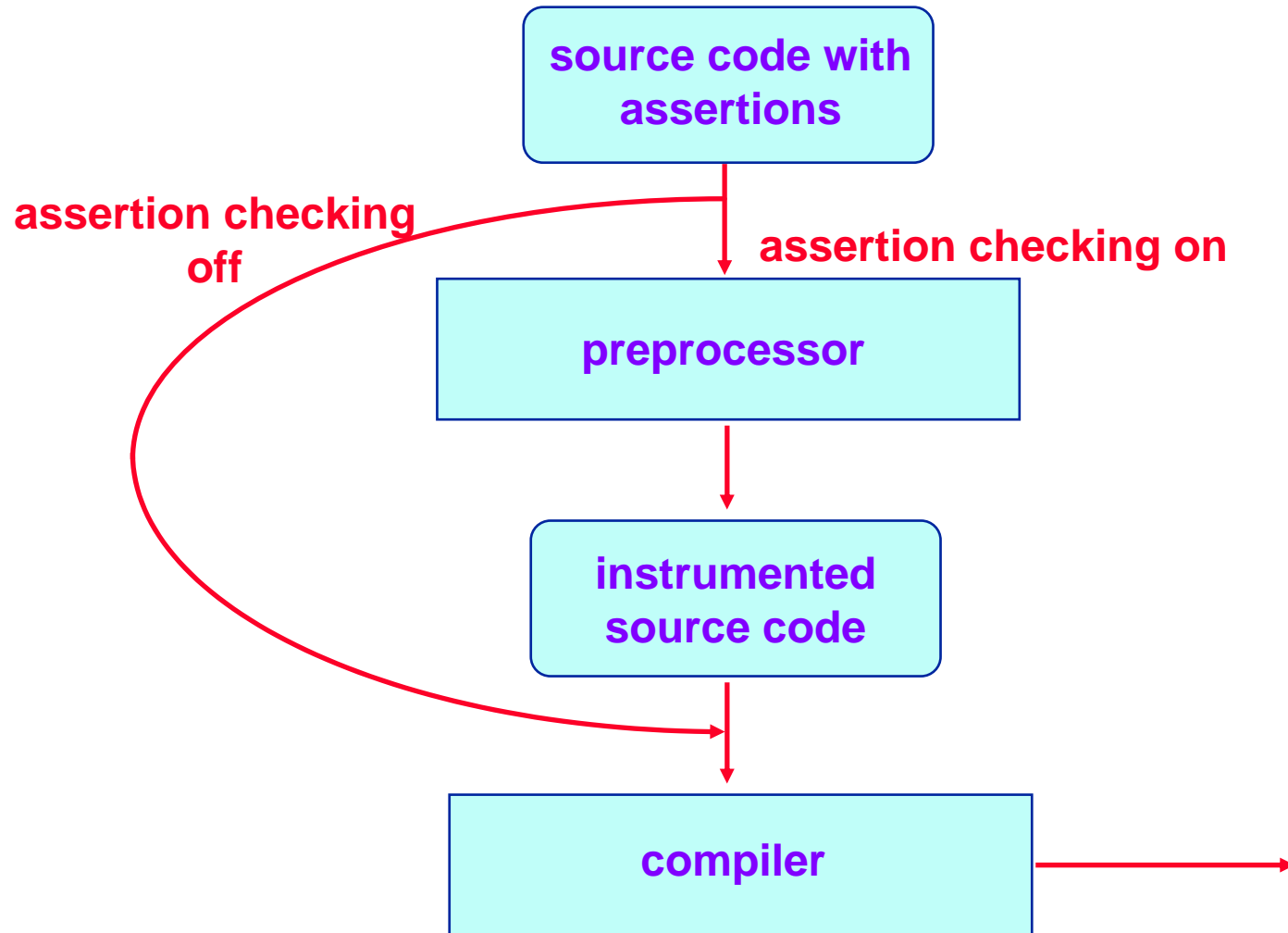
Parts of an Assertion Mechanism

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Execution Models

- **Suppress assertion checking**
 - Binary -> on or off
 - Multi-level
 - Select severity level to support
 - Suppress all assertions except those at severity level 3 and higher

Assertion preprocessor



Response model

- **Termination model**
 - When an assertion is violated, issue an error report and terminate
- **Failure and Warning model**
 - 2 (3) level model: failure, (warning,) no problem
 - On failure, issue an error report and terminate
 - On warning, issue an error report and continue
 - Continue as long as there is no problem

Annotation PreProcessor (APP)

- David Rosenblum
- APP supports 4 types of assertions
 - `assume--` specifies a precondition on a function
 - `promise--` specifies a postcondition on a function
 - `return--` specifies a constraint on the return value of a function
 - `assert--` specifies a constraint on an intermediate state of a function body
 - Where should these be placed?

APP

- provides quantification operations *all* and *some*
- default action gives an error message that prints information about the location of the violation and values of any variables involved
- users can define their own violation actions
- can associate severity levels with assertions
 - before processing indicate which severity levels should be checked

APP case study

- evaluated a program called Yeast
- 12,000 LOCs of C code
- half the program developed using APP
 - had specific rules for writing assertions
 - resulted in 116 assertions

Results of the study

- detected 19 known faults
 - 8 detected by APP
 - 6 would have been detected by APP, if it had been used
 - 2 detected by a dynamic storage certification routine
 - 3 required event sequence information
- Overhead of using assertions
 - 3.7% larger
 - no discernible difference in speed!

Example

```
public int binarySearch(int data [], int key){
    int lower = 0;
    int upper = data.length - 1;
    int location;
    while (true) {
        if(upper < lower)
            { return (-1) };
        else {
            location = midpoint(lower, upper);
            if (data [location] == key)
                {return (location); }
            else if (data[location] < key)
                {lower = location +1; }
            else
                { upper = location -1; }
        }
    }
}
```

Example: assume clause

```
public int binarySearch(int data [], int key){
    int lower = 0;
    int upper = data.length - 1;
    int location;
    while (true) {
        if(upper < lower)
            { return (-1) };
        else {
            location = midpoint(lower, upper);
            if (data [location] == key)
                {return (location); }
            else if (data[location] < key)
                {lower = location +1; }
            else
                { upper = location -1; }
        }
    }
}
```

assume

(data != null)&&

all (int i = 0; i < data.length - 1; i++) data[i] <= data[i + 1]

Example: return clause

```
public int binarySearch(int data [], int key){
    int lower = 0;
    int upper = data.length - 1;
    int location;
    while (true) {
        if(upper < lower)
            { return (-1) };
        else {
            location = midpoint(lower, upper);
            if (data [location] == key)
                {return (location); }
            else if (data[location] < key)
                {lower = location +1; }
            else
                { upper = location -1; }
        }
    }
}
```

return location where

```
((all(int i = 0; i < data.length; i++) data[i] != key) && (location == -1)) ||
((some(int i = 0; i < data.length; i++) data[i] == key) && (location == i))
```

Example: promise clause

```
public int binarySearch(int data [], int key){
    int lower = 0;
    int upper = data.length - 1;
    int location;
    while (true) {
        if(upper < lower)
            { return (-1) };
        else {
            location = midpoint(lower, upper);
            if (data [location] == key)
                {return (location); }
            else if (data[location] < key)
                {lower = location +1; }
            else
                { upper = location -1; }
        }
    }
}
```

`promise (data != null) &&`
`data.length == in data.length &&`
`all (int i = 0; i < data.length; i++) data[i] == in data[i]`

Example: internal assertions

```
public int binarySearch(int data [], int key){
    int lower = 0;
    int upper = data.length - 1;
    int location;
    while (true) {
        if(upper < lower)
            { return (-1) };
        else {
            location = midpoint(lower, upper);
            /** location is the midpoint between upper and lower
             * assert location <((float)(lower + upper)/2.0) + 1.0
             * assert location >((float)(lower + upper)/2.0) - 1.0
             */

            if (data [location] == key)
                {return (location); }
            else if (data[location] < key)
                {lower = location +1; }
            else
                { upper = location -1; }
        }
    }
}
```

Major objection to using assertions

- storage and runtime overhead
 - Not apparent in APP study
 - need more empirical data!!!
- optimization techniques could remove many of the assertions
 - basically proving that the assertion is valid
 - would expect that many of the assertions could be eliminated
 - preconditions are often redundant checks on the validity of the parameters

Assertions versus Exceptions

- **Assertion violation => error**
 - Predefined response
 - Error report
 - Terminate or continue
 - More expressive notation (e.g. All, Some, old, class invariant)
- **Exception violation => unusual case**
 - Style guideline=> exceptions should be reserved for truly exceptional situations
 - Program defined response
 - Handler
 - Different choices for resuming execution
 - Complex exception flow

Correct by Design

- Design/Code by contract
- Development method that incorporates assertions early in the design and coding process
- Eiffel, Bertrand Meyer included assertions as part of the language
- Assertions are the most requested feature in Java
 - 1.4 introduced a very limited assertion capability

Summary about Assertions

- assertions are a relatively easy way to improve software reliability
- assertion languages are accessible to most programmers
- assertions document intent and thus are useful beyond just runtime checking
- overhead is usually small, especially if optimization techniques are applied
- need more experimental data
 - which kinds of assertions are most useful?
 - what is the expected overhead?

Bottom line: Appears to be an effective approach with little overhead